

Energy-Aware Distributed Multimedia Systems

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Outline

- ◆ Introduction
- ◆ Concurrent Subband Modulation and Adaptive Decoding
- ◆ Energy-aware MPEG-4 FGS Video Streaming
- ◆ Power-aware Source Routing in MANETs
- ◆ Summary

Concurrent Subband Modulation and Adaptive Decoding

- ◆ Consider an ad hoc network of battery-powered mobile hosts
- ◆ Consider any point-to-point link between a sender and a receiver host
- ◆ Assume that in each time slot, the receiver provides an estimate of the channel characteristics to the transmitter as well as information about its remaining energy level
- ◆ Using this data and information about the remaining energy of the transmitter itself, the transmitter solves an optimization problem that would yield the energy optimal set of **modulation levels** and the **transmit power level**
- ◆ Given the set of transmitter parameters, the receiver must minimize its own energy consumption by selecting its Viterbi **decoder length**
- ◆ Constraints are the minimum throughput and the maximum bit-error-rate

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QAM and Adaptive Modulation

- ◆ In QAM, the bit error rate is strong function of the Signal-to-Noise ratio and the number of constellation points

$$BER = \frac{4}{b} \cdot \left(1 - \frac{1}{2^{b/2}}\right) \cdot Q\left(\sqrt{3 \cdot \frac{SNR}{2^b - 1}}\right)$$

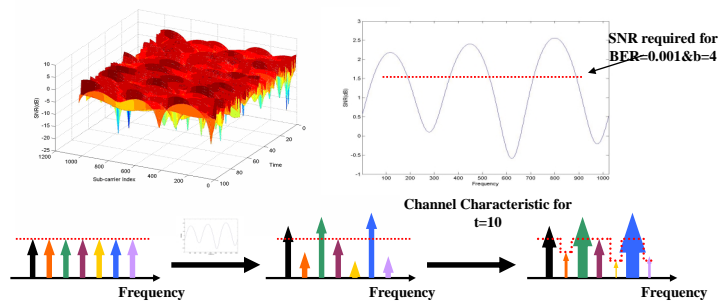
BER: Bit-Error-Rate

SNR: Received Signal-to-Noise Ratio

b : Number of bits per constellation point

Q : Normal cumulative probability function

- ◆ Key Idea : Choose the number of constellation points as a function of SNR in each sub-carrier



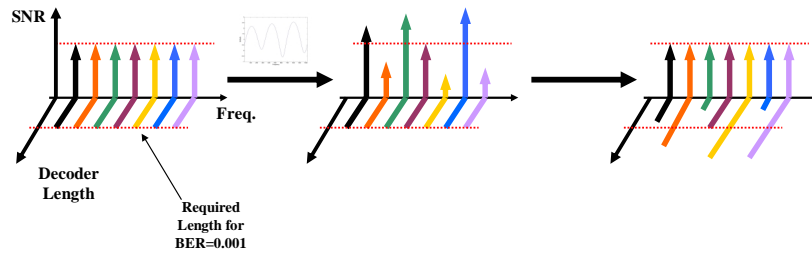
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Convolutional Codes and Adaptive Decoding

- ◆ Convolutional codes

- ◆ A Family of Forward Error Controlling codes
- ◆ Based on convolution of input data stream and two or more constant bit streams

- ◆ Key Idea : Adaptively change the decoding length based on received SNR for each sub-carrier



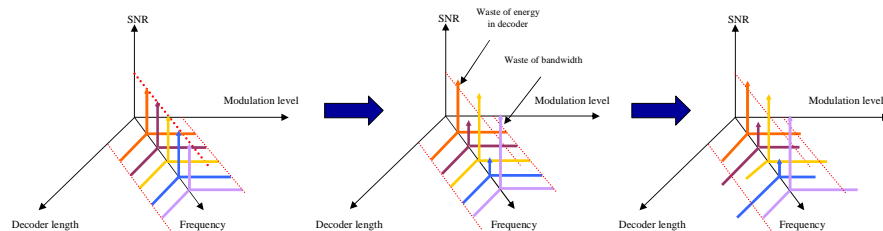
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Concurrent Adaptive Modulation and Decoding

- ◆ Using both adaptive modulation and decoding, i.e., having two degrees of freedom, the overall system energy can be reduced while keeping the average throughput constant

- ◆ Minimize the energy consumption per bit in the transmitter and the receiver of a wireless link while maintaining a minimum throughput and a maximum bit-error-rate

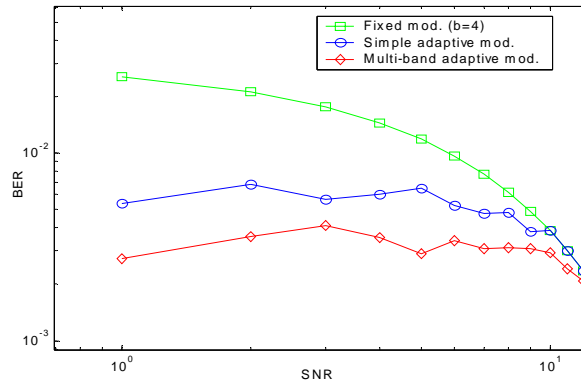
- ◆ Transmitter sets the set of modulation levels for all subbands and its transmit power level
- ◆ Receiver sets its Viterbi decoding length for each subband



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Results of Subband Modulation

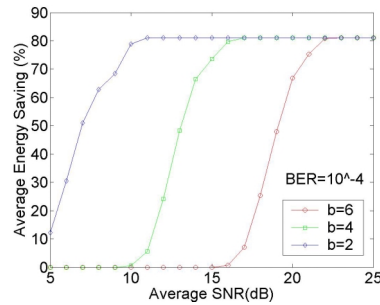
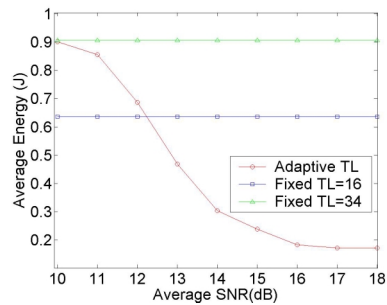
- A fixed decoder is used and the effect of subband modulation based on the estimated channel characteristic is reported



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Results of Adaptive Decoding

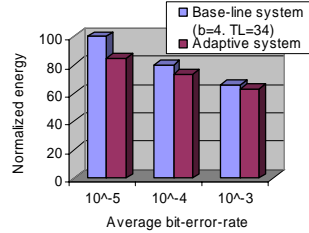
- A fixed modulation level is chosen and the effect of variable decoding length on average energy consumption is studied



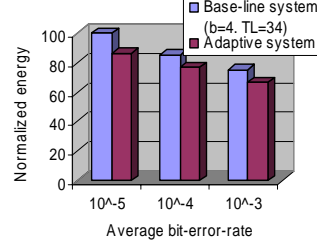
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Results: Total Energy Consumption

Overall Energy Saving (Alpha=0.7)

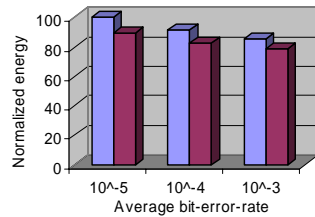


Overall Energy Saving (Alpha=0.5)



$$E_{avg} = \alpha \cdot E_{Transmit} + (1 - \alpha) \cdot E_{Receive}$$

Overall Energy Saving (Alpha=0.3)



- ◆ The base-line system is a system without any adaptive parameter except for the transmit power. Transmit power in base-line system is adaptively changed based on the average BER required on the receiver side. The optimized system refers to the adaptive transceiver proposed here

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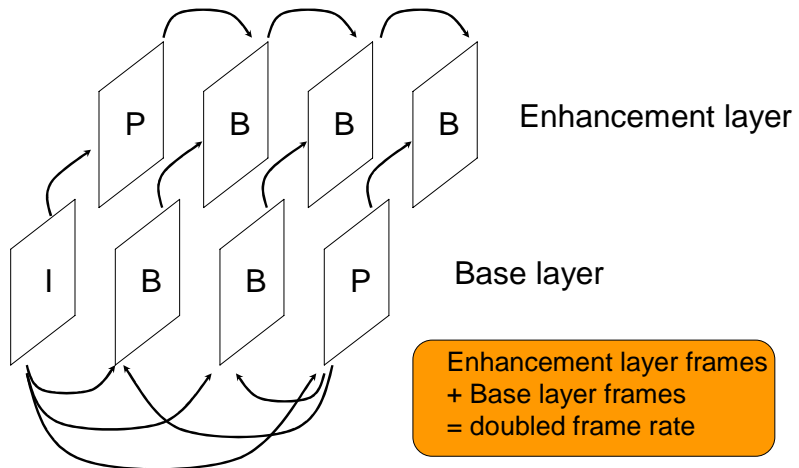
Wireless Video Streaming

- ◆ For mobile wireless video streaming two factors should be considered: high video quality and long service time
- ◆ Stable channel for real-time operation
 - ⊕ Video quality degradation due to channel congestion or error rate
 - ⊕ Scalable coding technique to be adaptive channel bandwidth variation
- ◆ Power-aware operation so as to increase the lifetime of battery-powered mobile system
 - ⊕ Optimal energy consumption to meet the required video quality
- ◆ Scalable video coding is required to be adaptive to channel variation
 - ⊕ a base layer (BL) + an enhancement layer (EL)

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Temporal Scalability

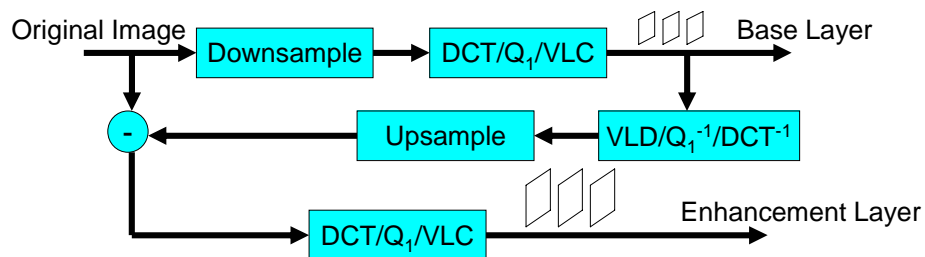
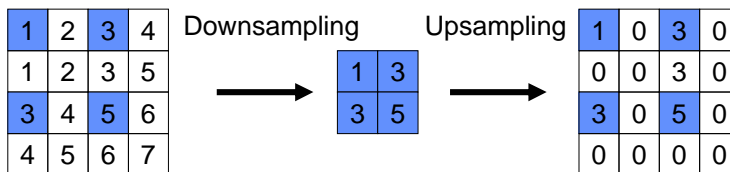
- Enhancement layer introduces a layer of P and B frames



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Spatial Scalability

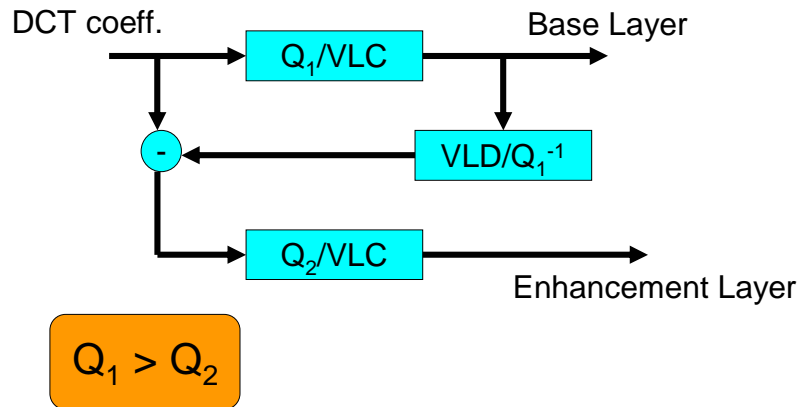
- Code a video into two layers at different spatial resolutions



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SNR Scalability

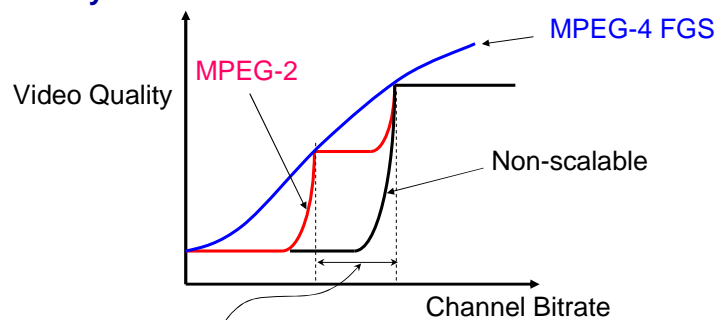
- ◆ Code a video into two layers with different quantization accuracies



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Deficiency of MPEG-2 Scalable Coding

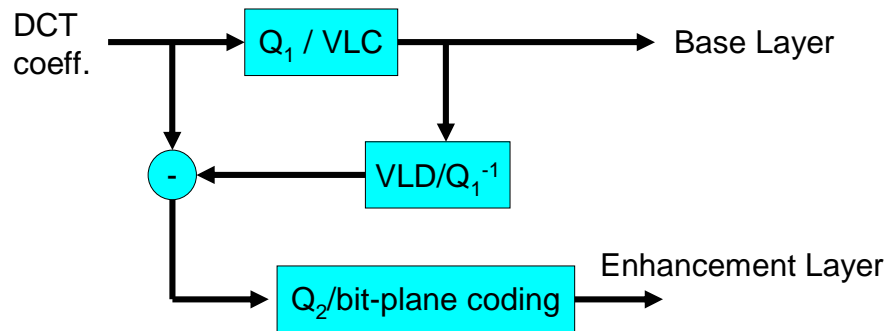
- ◆ MPEG-2 only provides two layers
 - ⊕ Dramatic change in the video quality as channel bandwidth varies
- ◆ MPEG-4 provides many more layers
 - ⊕ Continuous video quality improvement is desirable to maximally utilize current channel bandwidth



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MPEG-4 FGS Layer Structure

- ◆ Enhancement layer is equal to the original image minus the reconstructed image from the base layer



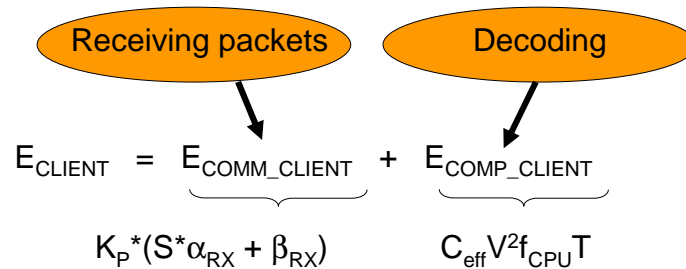
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Energy Consumption in Video Streaming

- ◆ There are two sources of energy consumption in wireless video streaming : the communication energy and the computation energy
 - ◆ Communication energy
 - ⊕ Transmitting packets (server)
 - ⊕ Receiving packets (client)
 - ◆ Computation energy
 - ⊕ Packetization (server)
 - ⊕ Decoding bit-streams (client)
- ◆ A video streaming system in which a server with infinite energy source and a battery-powered mobile client is considered

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Energy Consumption at the Client



K_p : number of packets
 S : packet size
 $\alpha_{\text{RX}}, \beta_{\text{RX}}$: regression coefficients

C_{eff} : effective capacitance
 V : operating voltage
 f_{CPU} : operating frequency
 T : streaming time

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Energy Waste at the Client

- ◆ **Video streaming is a real-time operation**
 - ✦ **If the client cannot process all the packets from the server in a given deadline, then the communication energy is wasted with no improvement of video quality**

Ex) Arrived packet count : A
Decoded packet count : M

Video quality = $\min(M, A)$

If $A > M$, then $(A-M)$ packets become useless, resulting in energy wasted in receiving these packets and ignoring them

- ◆ **For an energy-efficient streaming in which no energy waste occurs, A should be equal to M**

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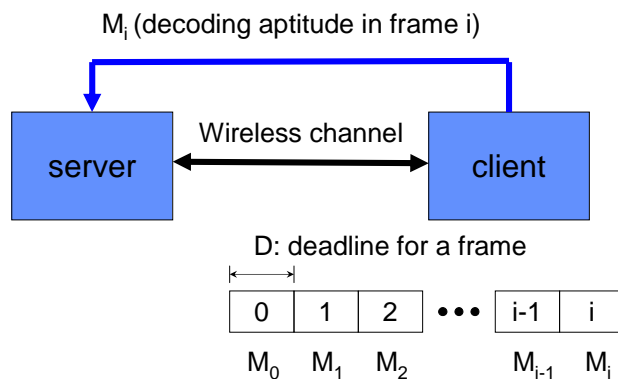
Decoding Aptitude of the Client

- ◆ **Decoding aptitude (M) of a mobile client**
 - ⊕ The amount of data that can be decoded in a given deadline
 - ⊕ M is proportional to the CPU frequency of the client
 - ⊕ M can be changed by dynamic voltage and frequency scaling (DVFS) or dynamic power management (DPM)
- ◆ N_i , normalized decoding load at time instance i , is defined as the ratio A_i/M_i
 - ⊕ A_i : the number of correctly arrived packets at the client
 - ⊕ M_i : decoding aptitude
- ◆ N_i represents the degree of energy waste and should be kept at one for no energy waste
- ◆ In order to make N_i one, the server should know M_i

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Client-feedback Video Streaming

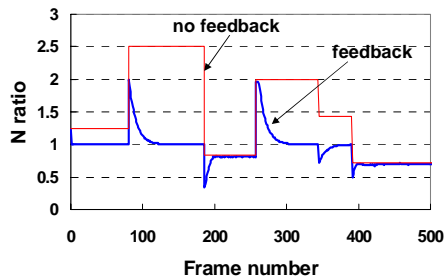
- ◆ The client sends a status packet to the server at regular time intervals
- ◆ The server sets the amount of data to be transferred based on the client status, making N_i equal to one



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Simulation Results on Apollo Testbed II

- ◆ Variations in N_i with different M_i/B ratios
 - ⊕ M_i : decoding aptitude at instance i
 - ⊕ B : maximum number of packets the server can send
- ◆ M_i trace : 0.8B, 0.4B, 1.2B, 0.5B, 0.7B, 1.4B
- ◆ Channel model : Gilbert-Elliot model with bit error rate (BER) of $1e-5$ and $1e-4$ for good and bad state, respectively



| | Energy waste | |
|------|--------------|-------|
| | No FB | FB |
| 0.8B | 18.74% | 0.21% |
| 0.4B | 57.35% | 2.57% |
| 1.2B | 0% | 0% |
| 0.5B | 48.49% | 6.27% |
| 0.7B | 28.69% | 0% |
| 1.4B | 0% | 0% |

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Experimental Setup

- ◆ Generated MPEG-4 FGS bit-streams using a QCIF test video

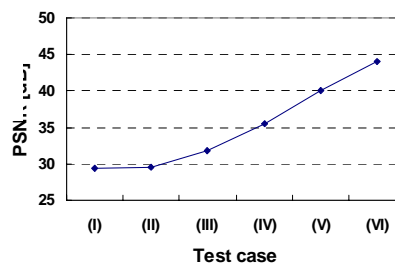
- ⊕ A base + FGS with five bit planes(bp0~4)
- ⊕ 256-byte packet size

- ◆ Six test cases

- ⊕ (I) base layer only
- ⊕ (II) base + bp0
- ⊕ (III) base + bp0 + bp1
- ⊕ (IV) base + bp0 + bp1 + bp2
- ⊕ (V) base + bp0 + bp1 + bp2 + bp3
- ⊕ (VI) base + bp0 + bp1 + bp2 + bp3 + bp4

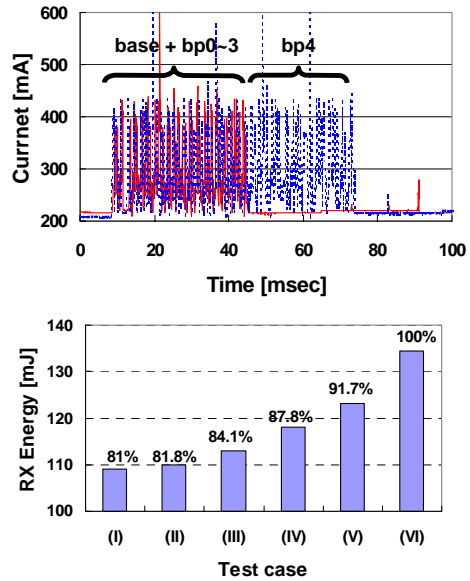
- ◆ Peak signal to noise ratio (PSNR) increases as more bit-planes are decoded

| | Size(byte) | Packet number |
|------------|------------|---------------|
| Base | 76 | 1 |
| FGS Header | 9 | 1 |
| bp0 | 18 | 1 |
| bp1 | 278 | 2 |
| bp2 | 1007 | 4 |
| bp3 | 2022 | 8 |
| bp4 | 3358 | 14 |



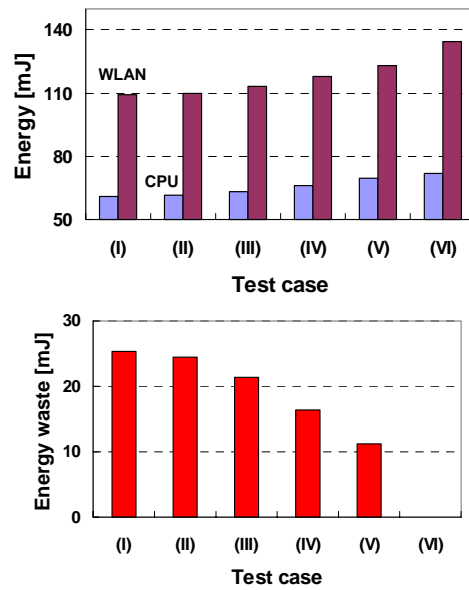
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Energy of Wireless LAN Card for Receiving Packets



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Energy Consumptions of the CPU and WLAN



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Energy-Aware Wireless Video Streaming in Adhoc Networks

- ◆ Given are a minimum throughput and a maximum bit-error-rate

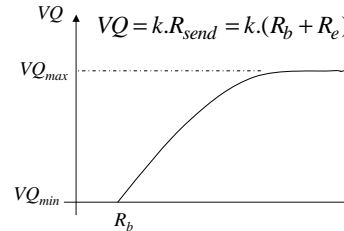
- ◆ Server's policy is to adopt
 - ⊕ Operating frequency of the CPU, f^S
 - ⊕ Transmit power, P_{amp}
 - ⊕ Set of modulation levels of the OFDM subcarriers, b_i

- ◆ Server's cost function is
 - ⊕ the total energy consumption of the client-server system in each frame

- ◆ Client's policy is to adopt the
 - ⊕ Operating frequency of CPU, f^C
 - ⊕ Truncation length of OFDM subcarriers, TL_i

- ◆ Client's cost function is
 - ⊕ the remaining lifetime difference between server and client

Fine Granularity Scalability feature of MPEG4



- ◆ The **Stackelberg game** theory is used to reduce the redundant energy consumption in a wireless video streaming system

- ◆ A maximum of 20% increase in the system lifetime is achieved. Moreover, average lifetime increases by 15%

- ◆ It is easy to transfer to the ADS bitsyX board

- ⊕ Needs protocol/Application modification

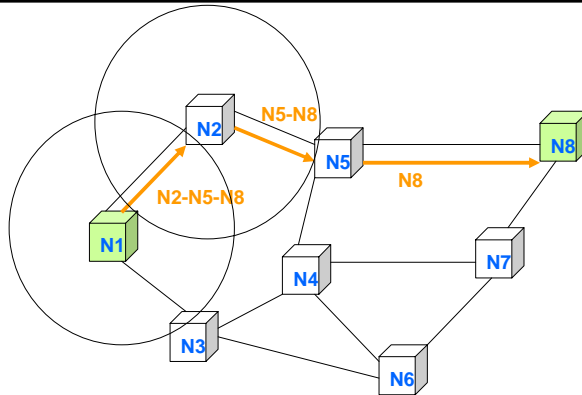
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Energy-Aware Routing in MANETs

- ◆ Power is one of the key design criteria in Wireless Adhoc Networks
- ◆ Network can become disconnected if some nodes die out due to lack of energy
- ◆ A design goal is to maximize the (useful) network lifetime by minimizing the variance of the remaining battery lifetime of the nodes in the network
- ◆ Our focus here is on on-demand (reactive) Multicast routing protocols
- ◆ Routing protocol ought to be responsive to the dynamics of the network (node density and mobility, energy discharge rates and traffic load)

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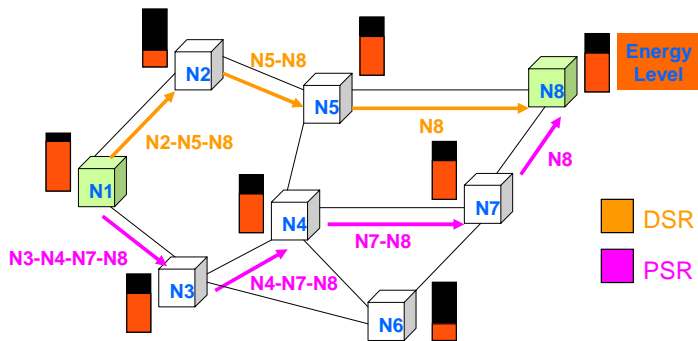
Dynamic Source Routing in MANET 101



- ◆ Assume fixed transmit power and fixed energy-conserving strategy in each node
- ◆ Route discovery is done by flooding the network with Route Requests
- ◆ Nodes promiscuously listen to control messages flowing through the network; Caching techniques improve performance considerably

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Power-aware Source Routing (PSR)



- ◆ Minimize sum of the energy cost of the links along the routing path
- ◆ Link cost is proportional to the inverse of the remaining battery capacity (residual energy) of the transmitting node

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PSR: Route Discovery

- ◆ **Similar to DSR, but with the following differences:**
 - ⊕ **An intermediate node passes on the first RREQ and all subsequent lower-cost RREQ's until a local timer expires**
 - ⊕ **Destination starts a timer after receiving the first RREQ and replies back only after that timer expires**
- ◆ **Control packet overhead for PSR is higher than that of DSR**

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PSR: Route Maintenance

- ◆ **Similar to DSR, but with the following differences :**
 - ⊕ **Path cost changes with time, and hence, cache entries of a node should not remain valid for a "long" period of time (aged cache entries should be purged from the cache tables of all nodes along the path)**
 - ⊕ **Each node in the path monitors the decrease in its residual energy from the time of route discovery; When this link cost increase goes beyond a threshold level, the node sends a route error back to the source as if the route was rendered invalid**
 - ⊕ **Invalidated routes are added to a victim buffer**
- ◆ **PSR results in a higher degree of energy balancing in the network**

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PSR: Exact Cost Function

$$\text{Min}_{\pi} C(\pi, t) = \sum_{i \in \pi} C_i(t)$$

where $C_i(t) = \rho_i \cdot \left(\frac{F_i}{E_{r,i}(t)} \right)^\alpha$

ρ_i : transmit power level of node i

F_i : full-charge battery capacity of node i

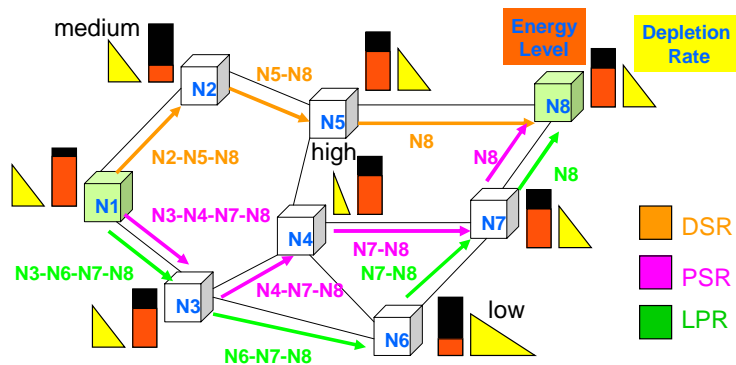
$E_{r,i}(t)$: remaining battery capacity of node i at time t

α : a positive weighting factor

- ◆ A cost is associated with every node on the path
- ◆ This cost is inversely proportional to the normalized residual energy of the node
- ◆ The cost function is graded, i.e., nodes with very low battery capacity dominate the total cost of the path

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Lifetime Prediction Routing (LPR)



- ◆ Maximize the minimum link cost along routing path
- ◆ Link cost is remaining lifetime of transmitting node
- ◆ Node lifetime is equal to the remaining battery capacity divided by the energy depletion rate

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LPR: Exact Cost Function

$$\text{Max}_{\pi} T_{\pi}(t) = \text{Min}_{i \in \pi} (T_i(t))$$

$T_{\pi}(t)$: lifetime of path π
 $T_i(t)$: predicted lifetime of node i in path π

$$T_i(t) = \frac{E_{r,i}(t)}{R_i(t)}$$

- ◆ $E_{r,i}(t)$: remaining energy of node i at time t
- ◆ $R_k(t)$: history-based rate of energy depletion of current node at time t
- ◆ N : length of the history used for calculating the simple moving average

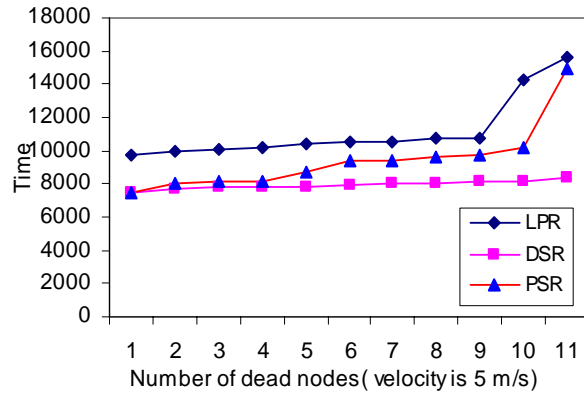
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Simulation Setup

- ◆ 20 nodes in 1000 x 1000 area with random way-point mobility
- ◆ 380 UDP connections which are randomly initiated between nodes at simulation time
- ◆ Key parameters of study are the network lifetime and RMS of energy consumption (Erms) in the network
 - ⊕ Effect of mobility
 - ⊕ Effect of radio transmission range

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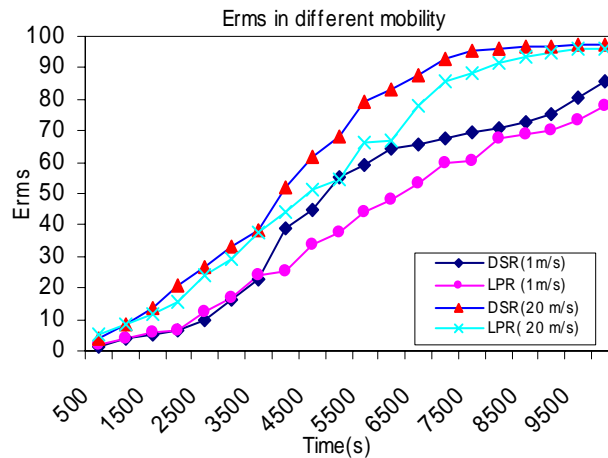
Network Lifetime



- ◆ LPR > PSR > DSR in terms of network lifetime
- ◆ Lifetime prediction is an effective technique to increase the network lifetime

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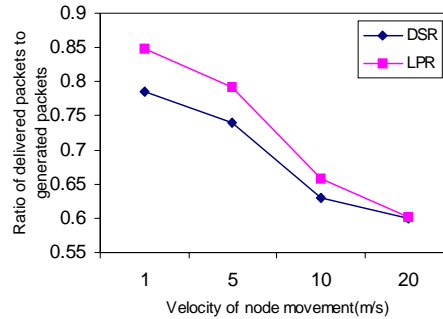
Effect of Mobility on Evaluation of Erms



- ◆ LPR is better than DSR in terms of energy variance

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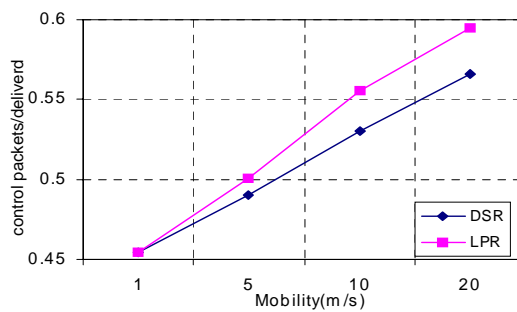
Effect of Mobility on Packet Delivery Ratio



- ◆ As velocity of node movement increases, ratio of packets delivered to packets generated for LPR goes down
- ◆ At higher velocities of node movement, LPR does not exhibit a significant gain over DSR

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Effect of Mobility on Control to Data Packet Ratio

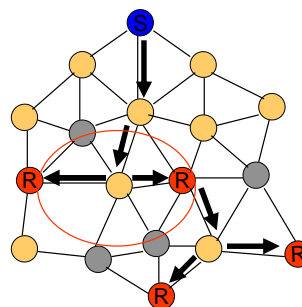


- ◆ Ratio of control packets to delivered packets in the network increases as velocity of node movement goes up for LPR
- ◆ Overhead of control packets is less than 6%

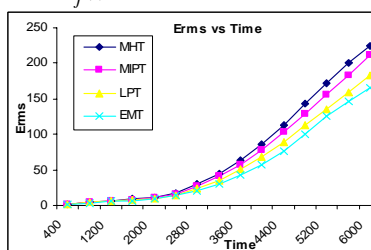
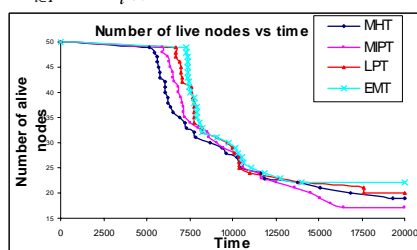
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Energy-aware On-demand Multicast Routing in Wireless Adhoc Networks

- ◆ Find a minimum-cost tree, which connects the source node S to all receivers in set R
- ◆ Maximize the network lifetime by minimizing the variance of the remaining battery lifetime of the nodes in the network
- ◆ There are two distinct methods to send out the packets from a multi-fanout node:
 - ◆ Multiple unicast
 - ◆ Single broadcast (more energy efficient)
- ◆ Must account for the “neighbor cost”



$$C(T, t) = \sum_{i \in T} \left\{ \rho_i \cdot \left(\frac{F_i}{R_i(t)} \right)^\alpha + \left(\text{if } \text{deg}(i) \geq 2 \text{ then } \sum_{j \in \text{neigh}(i) \wedge j \notin T} \left(\frac{F_j}{R_j(t)} \right)^\alpha \text{ else } 0 \right) \right\}$$



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Summary

- ◆ Energy consumption of an MPEG-2 decoder can be reduced by employing an application-aware DVFS strategy by as much as 20%
- ◆ By using a client-feedback method for the MPEG-4 streamer, about 20% reduction in the communication energy is achieved, which is up to 40% of the CPU energy
- ◆ Power-aware Routing protocols attempt to improve the network lifetime and balance both the remaining energy and energy depletion rate of the nodes in a MANET with low control overhead
- ◆ Performance of these protocols varies based on: Mobility, Radio transmission range, and Accuracy of lifetime prediction

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